Hamming It Up with Hamming Codes

CSE 461 Section Week 3

Error Detection/Correction

- We want to know when there are errors in communication
- Correcting errors would be even better! Why?
 - It'd save lots of time
- What are some ways we can correct errors?
 - Send data multiple times
 - Send longer symbols
 - Send data with the payload that's a function of the payload
 E.g., parity bits

Parity Bits

O Problem

• We want to send 1101

• Last bit gets flipped

• 1100 is sent instead

• How can we detect this?



• So instead of 1101 we send...

O 1101**1**

• If 11001 is received, we know that it's wrong—how?

• The parity bit for 1100 should be 0, but it's not... something was flipped!

• Even and odd parity



Error Detection/Correction

- Parity bits don't let us correct errors (by themselves)
- Can we do any better?
- What's the best way to detect and correct?



Hamming Codes: Background

- Richard Hamming
- Worked at Bell Labs
- Developed Hamming Codes to save time on punchcard reading errors



 Hamming codes now used for network communications as well as hard drive RAIDs



All other bits are message bits

00100101100

Bits in 1, 2, 4, 8, etc. positions are parity bits



- We want to send 1011100
- We put its message bits into the non-2ⁿ places, like so:



? ? 1 ? 0 1 1 ? 1 0 0 $P_1 p_2 m_3 p_4 m_5 m_6 m_7 p_8 m_9 m_{10} m_{11}$

- Each message bit is added to the parity bits that sum up to that message bit's place
- For m_3 , 3 = 2 + 1, so we add to p_1 and p_2

+1 +1 ? ? 1 ? 0 1 1 ? 1 0 0 $P_1 p_2 m_3 p_4 m_5 m_6 m_7 p_8 m_9 m_{10} m_{11}$



• For m_5 , 5 = 4 + 1, so we add to p_4 and p_1

• But $m_5 = 0$, so we don't add anything



• For m_6 , 6 = 4 + 2, so we add to p_4 and p_2





 For m₇, 7 = 4 + 2 + 1, so we add to p₄, p₂ and p₁

+1 +1 +1 +1 +1 +1 +1 ? ? 1 ? 0 1 1 ? 1 0 0 $P_1 p_2 m_3 p_4 m_5 m_6 m_7 p_8 m_9 m_{10} m_{11}$



• What would we do for m₉?

• 9 = 8 + 1, so we add to p_8 and p_1



• What would we do for m₁₀?

- 10 = 8 + 2, so we add to p_8 and p_2
- But $m_{10} = 0$, so we don't add anything



• What would we do for m₁₁?

- 11 = 8 + 2 + 1, so we add to p₈,
 p₂ and p₁
- But $m_{11} = 0$, so we don't add anything



Now we add up all of the parity bits

• What would each one be? (Even parity)



Now we add up all of the parity bits

• What would each one be? (Even parity)



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Now we add up all of the parity bits

• What would each one be? (Even parity)



Hamming Codes



- So we get 11100111100 as our bit string to send
- The receiver can recalculate the parity bits and make sure they match

Error Syndromes

The sender sent

11100111100

but what if we received

11000111100 ? Can we correct this?

Error Syndromes

Recalculate parity bits and you get the numbers in blue:

(We know there's an error because we didn't get 1101)

Error Syndromes

Add calculated parity bits to parity bits in received data:

1 1 0 0 0 1 1 1 1 0 0

 $P_1 p_2 m_3 p_4 m_5 m_6 m_7 p_8 m_9 m_{10} m_{11}$

Then reverse the sum and it will tell you the bit in error: 0011 -> third bit is wrong!

Now, you try! Decode this ASCII message (0b1000001 = 65 = 'A')

1011 0101 000

There may be bit errors! Assume each line encodes one byte of message data. (I.e., pad with a leading zero.)

1111 0011 001

1100 0001 111



Answers 1011 0101 000 1111 0011 001 -

1110 0001 111

Extract message bits

Convert to ASCII

1010 000 1001 001 1000 111



1100 0001 111

PIG